

Computational Intelligence

Winter Term 2009/10

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Design of Evolutionary Algorithms

Lecture 12

ad 2) design guidelines for variation operators

a) reachability

every $x \in X$ should be reachable from arbitrary $x_0 \in X$ after finite number of repeated variations with positive probability bounded from 0

b) unbiasedness

unless having gathered knowledge about problem variation operator should not favor particular subsets of solutions ⇒ formally: maximum entropy principle

c) control

variation operator should have parameters affecting shape of distributions; known from theory: weaken variation strength when approaching optimum

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Design of Evolutionary Algorithms

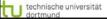
Lecture 12

Three tasks:

- 1. Choice of an appropriate problem representation.
- 2. Choice / design of variation operators acting in problem representation.
- 3. Choice of strategy parameters (includes initialization).

ad 1) different "schools":

- (a) operate on binary representation and define genotype/phenotype mapping
 - + can use standard algorithm
 - mapping may induce unintentional bias in search
- (b) no doctrine: use "most natural" representation
 - must design variation operators for specific representation
 - + if design done properly then no bias in search



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Design of Evolutionary Algorithms

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ad 2) design guidelines for variation operators in practice

binary search space $X = \mathbb{B}^n$

variation by k-point or uniform crossover and subsequent mutation

a) reachability:

regardless of the output of crossover we can move from $x \in \mathbb{B}^n$ to $y \in \mathbb{B}^n$ in 1 step with probability

$$p(x,y) = p_m^{H(x,y)} (1 - p_m)^{n-H(x,y)} > 0$$

where H(x,y) is Hamming distance between x and y.

Since min{ p(x,y): $x,y \in \mathbb{B}^n$ } = $\delta > 0$ we are done.

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b) unbiasedness

don't prefer any direction or subset of points without reason

⇒ use maximum entropy distribution for sampling!

properties:

- distributes probability mass as uniform as possible
- additional knowledge can be included as constraints:
- → under given constraints sample as uniform as possible

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Excursion: Maximum Entropy Distributions

Lecture 12

Knowledge available:

Discrete distribution with support $\{x_1, x_2, \dots x_n\}$ with $x_1 < x_2 < \dots x_n < \infty$

$$p_k = P\{X = x_k\}$$

⇒ leads to nonlinear constrained optimization problem:

$$-\sum_{k=1}^{n}p_{k}\log p_{k} \rightarrow \max!$$
 s.t.
$$\sum_{k=1}^{n}p_{k}=1$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p, a) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right)$$

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Formally:

Definition:

Let X be discrete random variable (r.v.) with $p_k = P\{X = x_k\}$ for some index set K. The quantity

 $H(X) = -\sum_{k \in K} p_k \log p_k$

is called the *entropy of the distribution* of X. If X is a continuous r.v. with p.d.f. $f_\chi(\cdot)$ then the entropy is given by

$$H(X) = \int_{-\infty}^{\infty} f_X(x) \log f_X(x) dx$$

The distribution of a random variable X for which H(X) is maximal is termed a maximum entropy distribution.

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Excursion: Maximum Entropy Distributions

Lecture 12

$$L(p,a) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right)$$

partial derivatives:

$$\begin{split} \frac{\partial L(p,a)}{\partial p_k} &= -1 - \log p_k + a \stackrel{!}{=} 0 \\ \frac{\partial L(p,a)}{\partial a} &= \sum_{k=1}^n p_k - 1 \stackrel{!}{=} 0 \\ \Rightarrow \sum_{k=1}^n p_k &= \sum_{k=1}^n e^{a-1} = n \, e^{a-1} \stackrel{!}{=} 1 \quad \Leftrightarrow \quad e^{a-1} = \frac{1}{n} \end{split} \qquad \begin{array}{c} p_k = \frac{1}{n} \\ \text{uniform distribution} \end{array}$$

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Excursion: Maximum Entropy Distributions

Lecture 12

Knowledge available:

Discrete distribution with support $\{1, 2, ..., n\}$ with $p_k = P\{X = k\}$ and E[X] = v

⇒ leads to nonlinear constrained optimization problem:

$$-\sum_{k=1}^n p_k \log p_k \quad \to \max!$$
 s.t.
$$\sum_{k=1}^n p_k = 1 \quad \text{ and } \quad \sum_{k=1}^n k \, p_k = \nu$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p, a, b) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right) + b \left(\sum_{k=1}^{n} k \cdot p_k - \nu\right)$$

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Excursion: Maximum Entropy Distributions

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$$\Rightarrow e^{a-1} = \frac{1}{\sum_{k=1}^{n} (c^b)^k} \Rightarrow p_k = e^{a-1+bk} = \frac{(e^b)^k}{\sum_{i=1}^{n} (e^b)^i}$$

discrete Boltzmann distribution $p_k = \frac{q^k}{\sum\limits_{i=1}^{n} q^i}$ $(q = e^b)$

value of g depends on v via third condition: (*)

$$\sum_{k=1}^{n} k p_{k} = \frac{\sum_{k=1}^{n} k q^{k}}{\sum_{i=1}^{n} q^{i}} = \frac{1 - (n+1) q^{n} + n q^{n+1}}{(1-q) (1-q^{n})} \stackrel{!}{=} \nu$$

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Excursion: Maximum Entropy Distributions

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$$L(p, a, b) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right) + b \left(\sum_{k=1}^{n} k \cdot p_k - \nu\right)$$

partial derivatives:

$$\frac{\partial L(p,a,b)}{\partial p_k} = -1 - \log p_k + a + b \, k \stackrel{!}{=} 0 \qquad \Rightarrow p_k = e^{a-1+b \, k}$$

$$\frac{\partial L(p,a,b)}{\partial a} = \sum_{k=1}^{n} p_k - 1 \stackrel{!}{=} 0$$

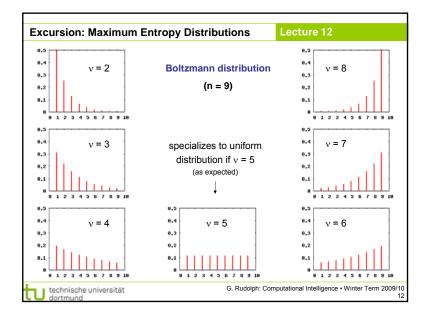
$$\frac{\partial L(p,a,b)}{\partial a} = \sum_{k=1}^{n} p_k - 1 \stackrel{!}{=} 0$$

$$\frac{\partial L(p,a,b)}{\partial b} \stackrel{(*)}{=} \sum_{k=1}^{n} k p_k - \nu \stackrel{!}{=} 0$$

$$\sum_{k=1}^{n} p_k = e^{a-1} \sum_{k=1}^{n} (e^b)^k \stackrel{!}{=} 1$$

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Excursion: Maximum Entropy Distributions

Lecture 12

Knowledge available:

Discrete distribution with support $\{1, 2, ..., n\}$ with E[X] = v and $V[X] = \eta^2$

⇒ leads to nonlinear constrained optimization problem:

$$-\sum_{k=1}^n p_k \log p_k \longrightarrow \max!$$

s.t.
$$\sum_{k=1}^n p_k = 1$$
 and $\sum_{k=1}^n k p_k = \nu$ and $\sum_{k=1}^n (k-\nu)^2 p_k = \eta^2$

solution: in principle, via Lagrange (find stationary point of Lagrangian function)

but very complicated analytically, if possible at all

⇒ consider special cases only

note: constraints are linear equations in p

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Excursion: Maximum Entropy Distributions

Lecture 12

Knowledge available:

Discrete distribution with unbounded support $\{0, 1, 2, ...\}$ and E[X] = v

⇒ leads to infinite-dimensional nonlinear constrained optimization problem:

$$-\sum_{k=0}^{\infty} p_k \log p_k \quad \to \max!$$

s.t.
$$\sum_{k=0}^{\infty} p_k = 1 \qquad \text{and} \qquad \sum_{k=0}^{\infty} k \, p_k = \nu$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p, a, b) = -\sum_{k=0}^{\infty} p_k \log p_k + a \left(\sum_{k=0}^{\infty} p_k - 1\right) + b \left(\sum_{k=0}^{\infty} k \cdot p_k - \nu\right)$$

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Excursion: Maximum Entropy Distributions

Lecture 12

Special case: n = 3 and E[X] = 2 and $V[X] = \eta^2$

Linear constraints uniquely determine distribution:

I.
$$p_1 + p_2 + p_3 = 1$$

$$n. p_1 + 2p_2 + 3p_3 = 2$$

II.
$$p_1 + p_2 + p_3 = 1$$
II. $p_1 + 2p_2 + 3p_3 = 2$
III. $p_1 + 0 + p_3 = \eta^2$

$$p_1 = \frac{\eta^2}{2}$$

II-I:
$$p_2 + 2p_3 = 1$$
 $p_3 = \frac{\eta^2}{2}$ $p_3 = \frac{\eta^2}{2}$

$$\Rightarrow p = \left(\frac{\eta^2}{2}, 1 - \eta^2, \frac{\eta^2}{2}\right) \quad \begin{cases} \eta^2 = \frac{1}{4} & \eta^2 = \frac{2}{3} \\ \frac{0.1}{0.2} & \frac{0.1}{0.2} \\ \frac{0.1}{0.2} & \frac{0.1}{0.2} \\ \frac{0.1}{0.2} & \frac{0.1}{0.2} \\ \frac{0.1}{0.2} & \frac{0.1}{0.2} \end{cases}$$

$$\eta^2 = \frac{1}{4}$$

$$\eta^2 = \frac{2}{3}$$

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Excursion: Maximum Entropy Distributions

$$L(p, a, b) = -\sum_{k=0}^{\infty} p_k \log p_k + a \left(\sum_{k=0}^{\infty} p_k - 1\right) + b \left(\sum_{k=0}^{\infty} k \cdot p_k - \nu\right)$$

partial derivatives:

$$\frac{\partial L(p,a,b)}{\partial p_k} = -1 - \log p_k + a + b \, k \stackrel{!}{=} 0 \qquad \Rightarrow p_k = e^{a-1+b \, k}$$

$$\frac{\partial L(p,a,b)}{\partial a} = \sum_{k=0}^{\infty} p_k - 1 \stackrel{!}{=} 0$$

$$\frac{\partial L(p,a,b)}{\partial b} \stackrel{\text{(*)}}{=} \sum_{k=0}^{\infty} k \, p_k - \nu \stackrel{!}{=} 0 \qquad \sum_{k=0}^{\infty} p_k = e^{a-1} \sum_{k=0}^{\infty} (e^b)^k \stackrel{!}{=} 1$$

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Excursion: Maximum Entropy Distributions

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$$\Rightarrow e^{a-1} = \frac{1}{\sum\limits_{k=0}^{\infty} (e^b)^k} \qquad \Rightarrow p_k = e^{a-1+bk} = \frac{(e^b)^k}{\sum\limits_{i=0}^{\infty} (e^b)^i}$$

set
$$q=e^b$$
 and insists that $q<1$ \Rightarrow $\sum\limits_{k=0}^{\infty}q^k=\frac{1}{1-q}$ insert

$$\Rightarrow p_k = (1-q)\,q^k \quad ext{for} \quad k=0,1,2,\dots \quad ext{geometrical distribution}$$

it remains to specify q; to proceed recall that
$$\sum_{k=0}^{\infty} k \, q^k = rac{q}{(1-q)^2}$$

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Lecture 12 **Excursion: Maximum Entropy Distributions** v = 1v = 7geometrical distribution 0.3 0.3 with E[x] = v-1 8 1 2 3 4 5 6 7 8 9 -1 8 1 2 3 4 5 6 7 8 9 v = 2v = 6p_k only shown for k = 0, 1, ..., 8v = 3v = 4v = 5G. Rudolph: Computational Intelligence • Winter Term 2009/10 technische universität

Excursion: Maximum Entropy Distributions

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⇒ value of q depends on v via third condition: (★)

$$\sum_{k=0}^{\infty} k \, p_k \, = \frac{\sum_{k=0}^{\infty} k \, q^k}{\sum_{i=0}^{\infty} q^i} \, = \, \frac{q}{1-q} \, \stackrel{!}{=} \, \nu$$

$$\Rightarrow q = \frac{\nu}{\nu + 1} = 1 - \frac{1}{\nu + 1}$$

$$\Rightarrow p_k = \frac{1}{\nu+1} \left(1 - \frac{1}{\nu+1} \right)^k$$

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