

Design of Evolutionary Algorithms	Lecture 10	Des	ign of E	Evolutio	onary	Algori	thms			Lectu	ire 10	
Genotype-Phenotype-Mapping $\mathbb{B}^n \to [L, R] \subset \mathbb{R}$ Genotype-Phenotype-Mapping $\mathbb{B}^n \to \mathbb{P}^{\log(n)}$ (example only)												
• Gray encoding for $b \in \mathbb{B}^n$ Let $a \in \mathbb{B}^n$ standard encoded. Then $b_i = \begin{cases} a_i, & \text{if } i = 1 \\ a_{i-1} \oplus a_i, & \text{if } i > 1 \end{cases} \oplus = XOR$			<ul> <li>e.g. standard encoding for b ∈ B<sup>n</sup></li> <li>individual:</li> </ul>									
000 001 011 010 110 111 101	100 - genotype		010	101	111	000	110	001	101	100	] ← genotyp	be
0 1 2 3 4 5 6	7 ← phenotype		0	1	2	3	4	5	6	7	← index	
OK, no hamming cliffs any longer $\Rightarrow$ small changes in phenotype "lead to" small chang since we consider evolution in terms of Darwin (not I $\Rightarrow$ small changes in genotype lead to small changes <b>but:</b> 1-Bit-change: 000 $\rightarrow$ 100 $\Rightarrow$ $\circledast$								-		ord / struct; permutation:	ex	
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Design of Evolutionary Algorithms         ad 1a) genotype-phenotype mapping         typically required: strong causality         → small changes in individual leads to small changes in f         → small changes in genotype should lead to small change	Design of Evolutionary Algorithms       Lecture 10         ad 1b) use "most natural" representation       typically required: strong causality         → small changes in individual leads to small changes in fitness       → need variation operators that obey that requirement											
but: how to find a genotype-phenotype mapping with that	but: how to find variation operators with that property?											
$\begin{tabular}{ c c c c }\hline \hline \textbf{necessary conditions}: \\ 1) g: \mathbb{B}^n \to X  can be computed efficiently (otherwise it is solution is solution of the second s$	otimal solution) ality endangered)	⇒r	eed des	ign guid	elines .							
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Design of Evolutionary Algorithms	Lecture 10	Design of Evolutionary Algorithms	Lecture 10						
ad 2) design guidelines for variation operators		ad 2) design guidelines for variation operators in practice							
<ul> <li>a) reachability <ul> <li>every x ∈ X should be reachable from arbitrary x<sub>0</sub> ∈ X</li> <li>after finite number of repeated variations with positive probability bounded from 0</li> </ul> </li> <li>b) unbiasedness <ul> <li>unless having gathered knowledge about problem</li> <li>variation operator should not favor particular subsets of solutions</li> <li>⇒ formally: maximum entropy principle</li> </ul> </li> <li>c) control <ul> <li>variation operator should have parameters affecting shape of distributions; known from theory: weaken variation strength when approaching optimum</li> </ul> </li> </ul>		binary search space $X = \mathbb{B}^n$ variation by k-point or uniform crossover and subsequent mutation a) <i>reachability</i> : regardless of the output of crossover we can move from $x \in \mathbb{B}^n$ to $y \in \mathbb{B}^n$ in 1 step with probability $p(x, y) = p_m^{H(x,y)} (1 - p_m)^{n - H(x,y)} > 0$ where $H(x,y)$ is Hamming distance between x and y. Since min{ $p(x,y): x, y \in \mathbb{B}^n$ } = $\delta > 0$ we are done.							
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b) <b>unbiasedness</b>		<u>Formally:</u>							
don't prefer any direction or subset of points without re	eason	Definition:							
$\Rightarrow$ use maximum entropy distribution for sampling!		Let X be discrete random variable (r.v.) with The quantity $H(X) = -\sum_{k \in K}$	with $p_k = P\{X = x_k\}$ for some index set K. $p_k \log p_k$						
properties: - distributes probability mass as uniform as possible - additional knowledge can be included as constrair → under given constraints sample as uniform as p	nts:	is called the <i>entropy of the distribution</i> $f_X(\cdot)$ then the entropy is given by $H(X) = -\int_{-\infty}^{\infty} f_X$ The distribution of a random variable X for <i>maximum entropy distribution</i> .	n of X. If X is a continuous r.v. with p.d.f. $_X(x) \log f_X(x) dx$						
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# Excursion: Maximum Entropy Distributions Lecture 10

## **Excursion: Maximum Entropy Distributions**

Lecture 10

Discrete distribution with support { 
$$x_1, x_2, ..., x_n$$
 } with  $x_1 < x_2 < ..., x_n < \infty$   
 $p_k = P\{X = x_k\}$ 

 $\Rightarrow$  leads to nonlinear constrained optimization problem:

$$-\sum_{k=1}^{n} p_k \log p_k \quad \to \max!$$
  
s.t. 
$$\sum_{k=1}^{n} p_k = 1$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p,a) = -\sum_{k=1}^{n} p_k \log p_k + a \left( \sum_{k=1}^{n} p_k - 1 \right)$$

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# Excursion: Maximum Entropy Distributions Lecture 10

### Knowledge available:

Discrete distribution with support { 1, 2, ..., n } with  $p_k = P \{X = k\}$  and E[X] = v

 $\Rightarrow$  leads to nonlinear constrained optimization problem:

$$-\sum_{k=1}^{n} p_k \log p_k \quad \rightarrow \max!$$
  
s.t. 
$$\sum_{k=1}^{n} p_k = 1 \quad \text{and} \quad \sum_{k=1}^{n} k p_k = \nu$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p, a, b) = -\sum_{k=1}^{n} p_k \log p_k + a \left( \sum_{k=1}^{n} p_k - 1 \right) + b \left( \sum_{k=1}^{n} k \cdot p_k - \nu \right)$$

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$$L(p,a) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right)$$

partial derivatives:

$$\frac{\partial L(p,a)}{\partial p_k} = -1 - \log p_k + a \stackrel{!}{=} 0 \qquad \Rightarrow p_k \stackrel{!}{=} e^{a-1}$$

$$\frac{\partial L(p,a)}{\partial a} = \sum_{k=1}^n p_k - 1 \stackrel{!}{=} 0 \qquad p_k = \frac{1}{n}$$

$$\Rightarrow \sum_{k=1}^n p_k = \sum_{k=1}^n e^{a-1} = n e^{a-1} \stackrel{!}{=} 1 \qquad \Leftrightarrow \qquad e^{a-1} = \frac{1}{n}$$

$$\lim_{k \to \infty} \frac{1}{n} e^{a-1} = \frac{1}{n}$$

## Excursion: Maximum Entropy Distributions

Lecture 10

$$L(p,a,b) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right) + b \left(\sum_{k=1}^{n} k \cdot p_k - \nu\right)$$

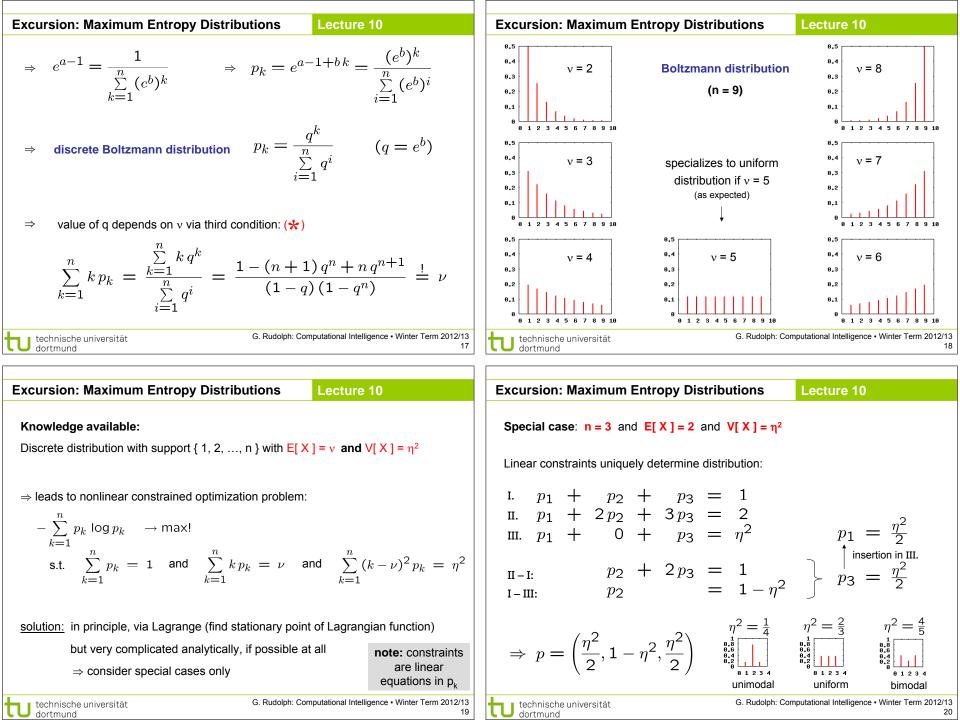
partial derivatives:

$$\frac{\partial L(p,a,b)}{\partial p_k} = -1 - \log p_k + a + b k \stackrel{!}{=} 0 \qquad \Rightarrow p_k = e^{a-1+b k}$$

$$\frac{\partial L(p,a,b)}{\partial a} = \sum_{k=1}^n p_k - 1 \stackrel{!}{=} 0$$

$$\frac{\partial L(p,a,b)}{\partial b} \stackrel{(\bigstar)}{=} \sum_{k=1}^n k p_k - \nu \stackrel{!}{=} 0 \qquad \sum_{k=1}^n p_k = e^{a-1} \sum_{k=1}^n (e^b)^k \stackrel{!}{=} 1$$

(continued on next slide)



#### **Excursion: Maximum Entropy Distributions** Lecture 10

Discrete distribution with unbounded support { 0, 1, 2, ... } and E[X] = v

 $\Rightarrow$  leads to infinite-dimensional nonlinear constrained optimization problem:

Knowledge available:

 $-\sum_{k=1}^{\infty} p_k \log p_k \longrightarrow \max!$ 

# **Excursion: Maximum Entropy Distributions**

Lecture 10

$$L(p,a,b) = -\sum_{k=0}^{\infty} p_k \log p_k + a \left(\sum_{k=0}^{\infty} p_k - 1\right) + b \left(\sum_{k=0}^{\infty} k \cdot p_k - \nu\right)$$

partial derivatives:

$$\frac{\partial L(p, a, b)}{\partial p_k} = -1 - \log p_k + a + b k \stackrel{!}{=} 0 \qquad \Rightarrow p_k = e^{a-1+b k}$$

$$\frac{\partial L(p, a, b)}{\partial a} = \sum_{k=0}^{\infty} p_k - 1 \stackrel{!}{=} 0$$

$$\frac{\partial L(p, a, b)}{\partial b} \stackrel{(\bigstar)}{=} \sum_{k=0}^{\infty} k p_k - \nu \stackrel{!}{=} 0 \qquad \sum_{k=0}^{\infty} p_k = e^{a-1} \sum_{k=0}^{\infty} (e^b)^k \stackrel{!}{=} 1$$
(continued on next slide)

s.t. 
$$\sum_{k=0}^{\infty} p_k = 1 \quad \text{and} \quad \sum_{k=0}^{\infty} k p_k = \nu$$
solution: via Lagrange (find stationary point of Lagrangian function)
$$L(p, a, b) = -\sum_{k=0}^{\infty} p_k \log p_k + a \left(\sum_{k=0}^{\infty} p_k - 1\right) + b \left(\sum_{k=0}^{\infty} k \cdot p_k - \nu\right)$$

$$\underbrace{\text{Lecture 10}}_{\text{Constrained}}$$
Excursion: Maximum Entropy Distributions
$$\underbrace{\text{Lecture 10}}_{k=0}$$

$$\Rightarrow \quad e^{a-1} = \frac{1}{\sum_{k=0}^{\infty} (e^b)^k} \quad \Rightarrow \quad p_k = e^{a-1+bk} = \frac{(e^b)^k}{\sum_{i=0}^{\infty} (e^b)^i}$$
set  $q = e^b$  and insists that  $q < 1 \quad \Rightarrow \quad \sum_{k=0}^{\infty} q^k = \frac{1}{1-q}$ 

$$\Rightarrow p_k = (1-q) q^k$$
 for  $k = 0, 1, 2, \dots$  geometrical distribution

it remains to specify q; to proceed recall that

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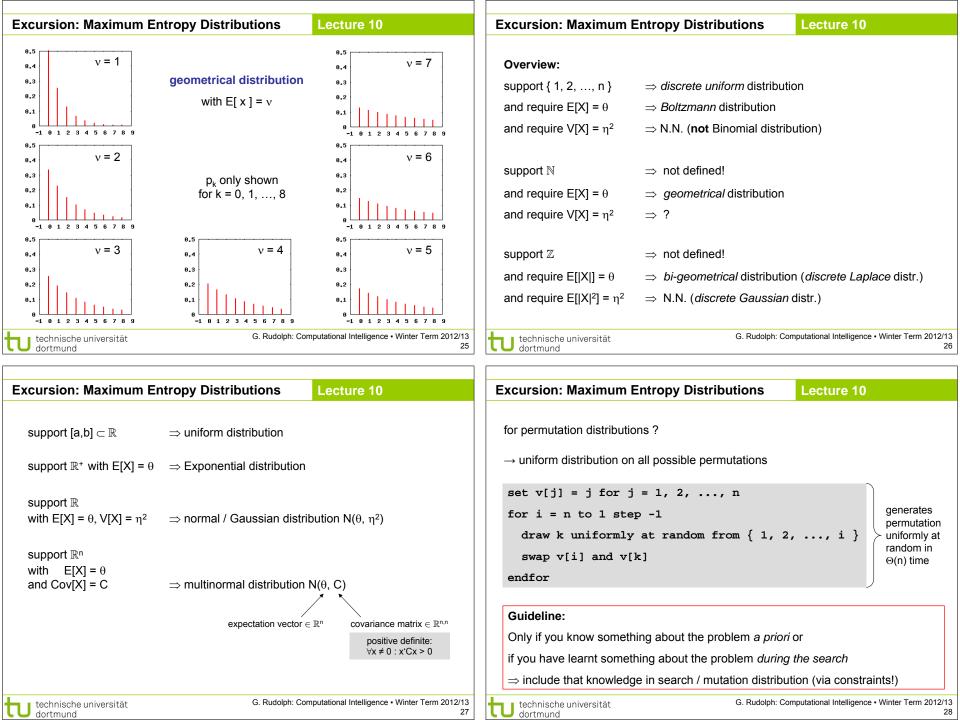
$$\sum_{k=0}^{\infty} k \, q^k \; = \; \frac{q}{(1-q)^2}$$

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**Excursion: Maximum Entropy Distributions** Lecture 10 value of q depends on v via third condition: (\*) $\Rightarrow$  $\sum_{k=0}^{\infty} k p_k = \frac{\sum_{k=0}^{\infty} k q^k}{\sum_{k=0}^{\infty} q^i} = \frac{q}{1-q} \stackrel{!}{=} \nu$  $\Rightarrow \quad q = \frac{\nu}{\nu+1} = 1 - \frac{1}{\nu+1}$  $\Rightarrow p_k = \frac{1}{\nu+1} \left( 1 - \frac{1}{\nu+1} \right)^k$ 



Design of Evolutionary Algorithms
 Lecture 10

 ad 2) design guidelines for variation operators in practice
 Image: Search space X = 2<sup>n</sup>

 a) reachability
 • every recombination results  
in some z = 2<sup>n</sup>

 a) reachability
 • every recombination results  
in some z = 2<sup>n</sup>

 b) unbiasedness
 • mean and the same metropy distribution of z = vary then teach  
probability in one step  
probability in one step  
probability on any the risk  
indary z = 2<sup>n</sup> with positive  
probability on any the risk  
of any z = 2<sup>n</sup> with positive  
probability of mutation should be 2<sup>n</sup>.

 ad b) need maximum entropy distribution over support 7<sup>n</sup>  
ad c) control variability by parameter  
-- formulate as constraint of maximum entropy distribution
 • React Compational functionary Algorithms
 Lecture 10

 Design of Evolutionary Algorithms
 Lecture 10

 result:  
a random variable Z with support 7, and probability distribution  
$$p_k := P\{Z = k\} = \frac{q}{2-q}(1-q)^{k}|$$
,  $k \in \mathbb{Z}$ ,  $q \in (0, 1)$   
symmetric with to 1, unimodal, spread manageable by q and has max entropy •  
generation of pseudo random numbers:
  $Z = G_1 - G_2$   
where  
 $U_i \sim U(0, 1) \Rightarrow G_i = \begin{bmatrix} log(1 - U_i) \\ log(1 - q_i) \end{bmatrix}$ ,  $i = 1, 2$ .  
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independent
  $0$ . React: Compatibulation long reactions of the reaction of the reaction of the reaction of the pendent
  $0$ . React: Compatibulation long reactions that the reaction of the reaction of pseudo random numbers:
  $Z = G_1 - G_2$   
where  
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