

Computational Intelligence

Winter Term 2020/21

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- Design of Evolutionary Algorithms
 - Design Guidelines
 - Genotype-Phenotype Mapping
 - Maximum Entropy Distributions

Three tasks:

- 1. Choice of an appropriate problem representation.
- 2. Choice / design of variation operators acting in problem representation.
- 3. Choice of strategy parameters (includes initialization).

- ad 1) different "schools":
 - (a) operate on binary representation and define genotype/phenotype mapping
 - + can use standard algorithm
 - mapping may induce unintentional bias in search

(b) no doctrine: use "most natural" representation

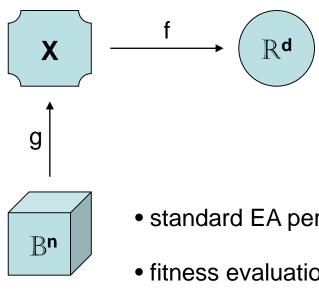
- must design variation operators for specific representation
- + if design done properly then no bias in search

Design of Evolutionary Algorithms

ad 1a) genotype-phenotype mapping

original problem f: $X \to \mathbb{R}^d$

scenario: no standard algorithm for search space X available



- \bullet standard EA performs variation on binary strings $b \in \mathbb{B}^n$
- fitness evaluation of individual b via (f ∘ g)(b) = f(g(b))
 where g: Bⁿ → X is genotype-phenotype mapping
- selection operation independent from representation

4

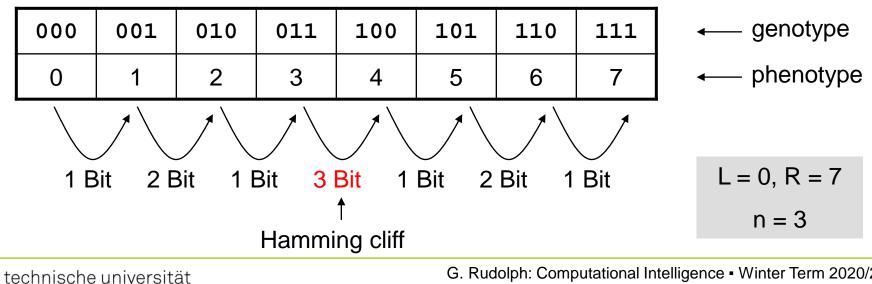
Genotype-Phenotype-Mapping $\mathbb{B}^n \to [L, R] \subset \mathbb{R}$

• Standard encoding for $b \in \mathbb{B}^n$

$$x = L + \frac{R - L}{2^n - 1} \sum_{i=0}^{n-1} b_{n-i} 2^i$$

 \rightarrow Problem: hamming cliffs

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Genotype-Phenotype-Mapping $\mathbb{B}^n \to [L, R] \subset \mathbb{R}$

 \bullet Gray encoding for $b \in \mathbb{B}^n$

Let $a \in \mathbb{B}^n$ standard encoded. Then $b_i = \begin{cases} a_i, & \text{if } i = 1 \\ a_{i-1} \oplus a_i, & \text{if } i > 1 \end{cases} \oplus = XOR$

000	001	011	010	110	111	101	100	← genotype
0	1	2	3	4	5	6	7	phenotype

OK, no hamming cliffs any longer ...

 \Rightarrow small changes in phenotype "lead to" small changes in genotype

since we consider evolution in terms of Darwin (not Lamarck):

 \Rightarrow small changes in genotype lead to small changes in phenotype!

but: 1-Bit-change: $000 \rightarrow 100 \Rightarrow \bigotimes$

Genotype-Phenotype-Mapping $\mathbb{B}^n \to \mathbb{P}^{\log(n)}$ (example only)

 \bullet e.g. standard encoding for $b \in \mathbb{B}^n$

individual:

010	101	111	000	110	001	101	100	← genotype
0	1	2	3	4	5	6	7	← index

consider index and associated genotype entry as unit / record / struct;

sort units with respect to genotype value, old indices yield permutation:

000	001	010	100	101	101	110	111	← genotype
3	5	0	7	1	6	4	2	← old index

= permutation

7

ad 1a) genotype-phenotype mapping

typically required: strong causality

- \rightarrow small changes in individual leads to small changes in fitness
- \rightarrow small changes in genotype should lead to small changes in phenotype

but: how to find a genotype-phenotype mapping with that property?

necessary conditions:

- 1) g: $\mathbb{B}^n \to X$ can be computed efficiently (otherwise it is senseless)
- 2) g: $\mathbb{B}^n \to X$ is surjective (otherwise we might miss the optimal solution)
- 3) g: $\mathbb{B}^n \to X$ preserves closeness (otherwise strong causality endangered)

Let d(\cdot , \cdot) be a metric on \mathbb{B}^n and d_X(\cdot , \cdot) be a metric on X.

 $\forall x, \, y, \, z \, \in \, \mathbb{B}^n \colon d(x, \, y) \leq d(x, \, z) \, \Rightarrow d_X(g(x), \, g(y)) \leq d_X(g(x), \, g(z))$

ad 1b) use "most natural" representation

typically required: strong causality

- \rightarrow small changes in individual leads to small changes in fitness
- \rightarrow need variation operators that obey that requirement

but: how to find variation operators with that property?

 \Rightarrow need design guidelines ...



ad 2) design guidelines for variation operators

a) reachability

every $x \in X$ should be reachable from arbitrary $x_0 \in X$ after finite number of repeated variations with positive probability bounded from 0

b) unbiasedness

unless having gathered knowledge about problem variation operator should not favor particular subsets of solutions \Rightarrow formally: <u>maximum entropy principle</u>

c) control

variation operator should have parameters affecting shape of distributions; known from theory: weaken variation strength when approaching optimum



ad 2) design guidelines for variation operators in practice

binary search space $X = \mathbb{B}^n$

variation by k-point or uniform crossover and subsequent mutation

a) *reachability*:

regardless of the output of crossover we can move from $x \in \mathbb{B}^n$ to $y \in \mathbb{B}^n$ in 1 step with probability

$$p(x,y) = p_m^{H(x,y)} (1 - p_m)^{n - H(x,y)} > 0$$

where H(x,y) is Hamming distance between x and y.

Since $\min\{p(x,y): x, y \in \mathbb{B}^n\} = \delta > 0$ we are done.

b) **unbiasedness**

don't prefer any direction or subset of points without reason

 \Rightarrow use maximum entropy distribution for sampling!

properties:

- distributes probability mass as uniform as possible
- additional knowledge can be included as constraints:
 → under given constraints sample as uniform as possible



Formally:

Definition:

Let X be discrete random variable (r.v.) with $p_k = P\{X = x_k\}$ for some index set K. The quantity

$$H(X) = -\sum_{k \in K} p_k \log p_k$$

is called the *entropy of the distribution* of X. If X is a continuous r.v. with p.d.f. $f_X(\cdot)$ then the entropy is given by

$$H(X) = -\int_{-\infty}^{\infty} f_X(x) \log f_X(x) dx$$

The distribution of a random variable X for which H(X) is maximal is termed a *maximum entropy distribution*.



Knowledge available:

Discrete distribution with support { $x_1, x_2, ..., x_n$ } with $x_1 < x_2 < ..., x_n < \infty$

$$p_k = \mathsf{P}\{X = x_k\}$$

Lecture 06

 \Rightarrow leads to nonlinear constrained optimization problem:

$$-\sum_{k=1}^{n} p_k \log p_k \rightarrow \max!$$

s.t.
$$\sum_{k=1}^{n} p_k = 1$$

solution: via Lagrange (find stationary point of Lagrangian function)

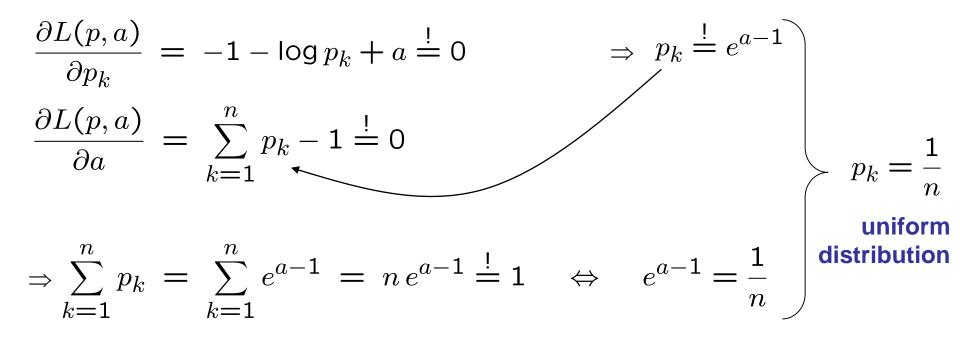
$$L(p,a) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1 \right)$$

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$$L(p,a) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1 \right)$$

partial derivatives:



Knowledge available:

 \mathbf{n}

Discrete distribution with support { 1, 2, ..., n } with $p_k = P \{ X = k \}$ and E[X] = v

 \Rightarrow leads to nonlinear constrained optimization problem:

$$-\sum_{k=1}^{n} p_k \log p_k \rightarrow \max!$$

s.t.
$$\sum_{k=1}^{n} p_k = 1 \text{ and } \sum_{k=1}^{n} k p_k = \nu$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p, a, b) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1 \right) + b \left(\sum_{k=1}^{n} k \cdot p_k - \nu \right)$$

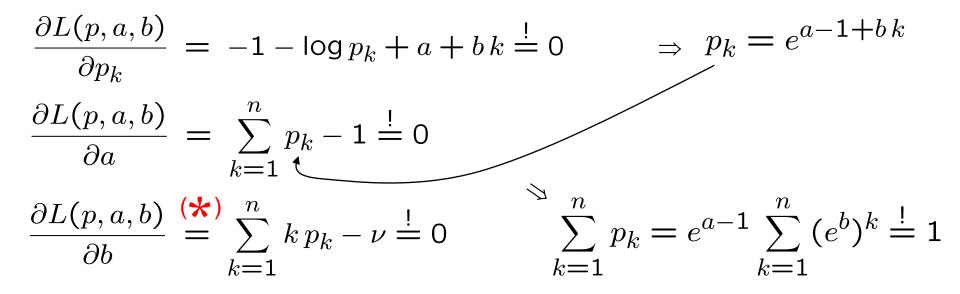
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Lecture 06

Lecture 06

$$L(p,a,b) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right) + b \left(\sum_{k=1}^{n} k \cdot p_k - \nu\right)$$

partial derivatives:



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Lecture 06

$$\Rightarrow e^{a-1} = \frac{1}{\sum_{k=1}^{n} (e^b)^k} \qquad \Rightarrow p_k = e^{a-1+bk} = \frac{(e^b)^k}{\sum_{i=1}^{n} (e^b)^i}$$

⇒ discrete Boltzmann distribution

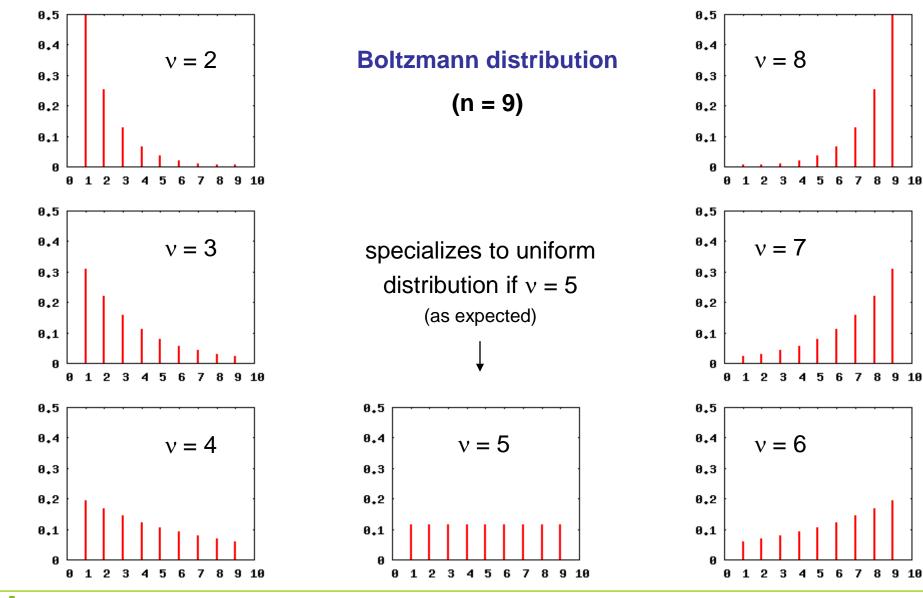
$$p_k = \frac{q^k}{\sum\limits_{i=1}^n q^i} \qquad (q = e^b)$$

 \Rightarrow value of q depends on v via third condition: (\bigstar)

$$\sum_{k=1}^{n} k p_k = \frac{\sum_{k=1}^{n} k q^k}{\sum_{i=1}^{n} q^i} = \frac{1 - (n+1) q^n + n q^{n+1}}{(1-q) (1-q^n)} \stackrel{!}{=} \nu$$

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Lecture 06



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Knowledge available:

Discrete distribution with support { 1, 2, ..., n } with E[X] = v and V[X] = η^2

 \Rightarrow leads to nonlinear constrained optimization problem:

$$-\sum_{k=1}^{n} p_k \log p_k \longrightarrow \max!$$

s.t.
$$\sum_{k=1}^{n} p_k = 1 \quad \text{and} \quad \sum_{k=1}^{n} k p_k = \nu \quad \text{and} \quad \sum_{k=1}^{n} (k-\nu)^2 p_k = \eta^2$$

solution:in principle, via Lagrange (find stationary point of Lagrangian function)but very complicated analytically, if possible at allnote: constraints
are linear
equations in p_k

Lecture 06

Special case: n = 3 and E[X] = 2 and $V[X] = \eta^2$

Linear constraints uniquely determine distribution:

I.
$$p_1 + p_2 + p_3 = 1$$

II. $p_1 + 2p_2 + 3p_3 = 2$
III. $p_1 + 0 + p_3 = \eta^2$
II. $p_1 + 0 + p_3 = \eta^2$
II. $p_2 + 2p_3 = 1$
 $p_3 = \frac{\eta^2}{2}$
 $p_3 = \frac{\eta^2}{2}$

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Knowledge available:

Discrete distribution with unbounded support { 0, 1, 2, ... } and E[X] = v

 \Rightarrow leads to <u>infinite-dimensional</u> nonlinear constrained optimization problem:

$$-\sum_{k=0}^{\infty} p_k \log p_k \to \max!$$

s.t.
$$\sum_{k=0}^{\infty} p_k = 1 \quad \text{and} \quad \sum_{k=0}^{\infty} k p_k = \nu$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p, a, b) = -\sum_{k=0}^{\infty} p_k \log p_k + a \left(\sum_{k=0}^{\infty} p_k - 1\right) + b \left(\sum_{k=0}^{\infty} k \cdot p_k - \nu\right)$$

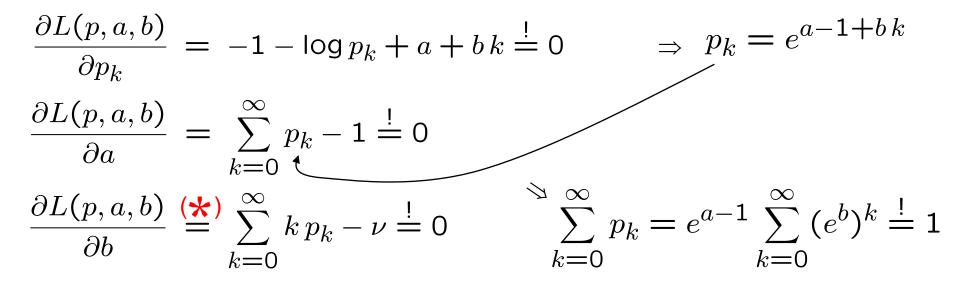
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$$L(p,a,b) = -\sum_{k=0}^{\infty} p_k \log p_k + a \left(\sum_{k=0}^{\infty} p_k - 1\right) + b \left(\sum_{k=0}^{\infty} k \cdot p_k - \nu\right)$$

partial derivatives:



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Lecture 06

$$\Rightarrow e^{a-1} = \frac{1}{\sum_{k=0}^{\infty} (e^b)^k} \Rightarrow p_k = e^{a-1+bk} = \frac{(e^b)^k}{\sum_{i=0}^{\infty} (e^b)^i}$$

set $q = e^b$ and insists that $q < 1 \Rightarrow \sum_{k=0}^{\infty} q^k = \frac{1}{1-q}$ insert

 $\Rightarrow p_k = (1 - q) q^k$ for k = 0, 1, 2, ... geometrical distribution

it remains to specify q; to proceed recall that

$$\sum_{k=0}^{\infty} k \, q^k \; = \; \frac{q}{(1-q)^2}$$



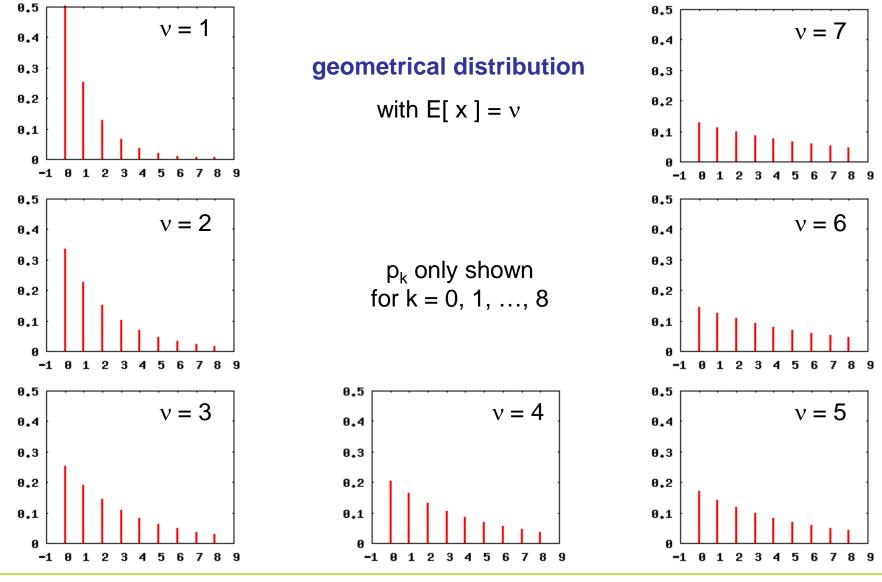
- Lecture 06
- \Rightarrow value of q depends on v via third condition: (*)

$$\sum_{k=0}^{\infty} k p_k = \frac{\sum_{k=0}^{\infty} k q^k}{\sum_{i=0}^{\infty} q^i} = \frac{q}{1-q} \stackrel{!}{=} \nu$$

$$\Rightarrow \quad q = \frac{\nu}{\nu+1} = 1 - \frac{1}{\nu+1}$$

$$\Rightarrow p_k = \frac{1}{\nu+1} \left(1 - \frac{1}{\nu+1} \right)^k$$

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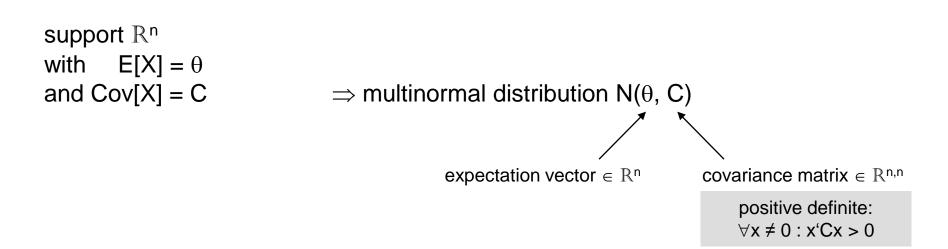
Overview:

support { 1, 2,, n }	\Rightarrow discrete uniform distribution
and require $E[X] = \theta$	\Rightarrow <i>Boltzmann</i> distribution
and require V[X] = η^2	\Rightarrow N.N. (not Binomial distribution)
support \mathbb{N}	\Rightarrow not defined!
and require $E[X] = \theta$	\Rightarrow geometrical distribution
and require V[X] = η^2	\Rightarrow ?
support \mathbb{Z}	\Rightarrow not defined!
and require $E[X] = \theta$	\Rightarrow <i>bi-geometrical</i> distribution (<i>discrete Laplace</i> distr.)
and require $E[X ^2] = \eta^2$	⇒ N.N. (<i>discrete Gaussian</i> distr.)

support [a,b] $\subset \mathbb{R}$ \Rightarrow uniform distribution

support \mathbb{R}^+ with $E[X] = \theta \implies$ Exponential distribution

support R with E[X] = θ , V[X] = $\eta^2 \implies$ normal / Gaussian distribution N(θ , η^2)





Lecture 06

for permutation distributions ?

 \rightarrow uniform distribution on all possible permutations

```
set v[j] = j for j = 1, 2, ..., n
for i = n to 1 step -1
    draw k uniformly at random from { 1, 2, ..., i }
    swap v[i] and v[k]
endfor
```

Guideline:

Only if you know something about the problem a priori or

if you have learnt something about the problem *during the search*

 \Rightarrow include that knowledge in search / mutation distribution (via constraints!)