

# **Computational Intelligence**

Winter Term 2021/22

Prof. Dr. Günter Rudolph

Lehrstuhl für Algorithm Engineering (LS 11)

Fakultät für Informatik

TU Dortmund

technische universität

**Plan for Today** 

G. Rudolph: Computational Intelligence • Winter Term 2021/22

Lecture 06

#### **Design of Evolutionary Algorithms**

Lecture 06

#### Three tasks:

- 1. Choice of an appropriate problem representation.
- 2. Choice / design of variation operators acting in problem representation.
- 3. Choice of strategy parameters (includes initialization).

ad 1) different "schools":

- (a) operate on binary representation and define genotype/phenotype mapping
  - + can use standard algorithm
  - mapping may induce unintentional bias in search
- (b) no doctrine: use "most natural" representation
  - must design variation operators for specific representation
  - + if design done properly then no bias in search

Lecture 06

ad 1a) genotype-phenotype mapping

**Design of Evolutionary Algorithms** 

• Design of Evolutionary Algorithms

**Design Guidelines** 

Genotype-Phenotype Mapping

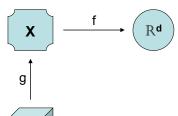
**Maximum Entropy Distributions** 

original problem  $f: X \to \mathbb{R}^d$ 

 $\mathbb{B}^{\mathsf{n}}$ 

■ technische universität dortmund

scenario: no standard algorithm for search space X available



- standard EA performs variation on binary strings b  $\in \mathbb{B}^n$
- fitness evaluation of individual b via (f g)(b) = f(g(b)) where g:  $\mathbb{B}^n \to X$  is genotype-phenotype mapping
- selection operation independent from representation

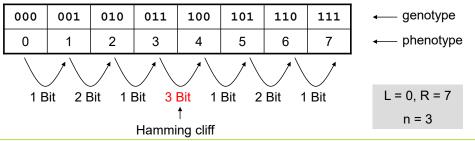
#### **Design of Evolutionary Algorithms**

Lecture 06

#### **Genotype-Phenotype-Mapping** $B^n \rightarrow [L, R] \subset R$

$$x = L + \frac{R - L}{2^{n} - 1} \sum_{i=0}^{n-1} b_{n-i} 2^{i}$$

→ Problem: *hamming cliffs* 



technische universität dortmund

G. Rudolph: Computational Intelligence • Winter Term 2021/22

# ullet Standard encoding for $b \in \mathbb{B}^n$

#### **Genotype-Phenotype-Mapping** $\mathbb{B}^n \to [L, R] \subset \mathbb{R}$

• Gray encoding for  $b \in B^n$ 

**Design of Evolutionary Algorithms** 

Let 
$$a \in \mathbb{B}^n$$
 standard encoded. Then  $b_i = \begin{cases} a_i, & \text{if } i = 1 \\ a_{i-1} \oplus a_i, & \text{if } i > 1 \end{cases}$ 

000	001	011	010	110	111	101	100	←— genotype
0	1	2	3	4	5	6	7	← phenotype

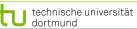
OK, no hamming cliffs any longer ...

⇒ small changes in phenotype "lead to" small changes in genotype

since we consider evolution in terms of Darwin (not Lamarck):

⇒ small changes in genotype lead to small changes in phenotype!

**but:** 1-Bit-change:  $000 \rightarrow 100 \Rightarrow \otimes$ 



G. Rudolph: Computational Intelligence • Winter Term 2021/22

Lecture 06

#### **Design of Evolutionary Algorithms**

Lecture 06

**Genotype-Phenotype-Mapping**  $\mathbb{B}^n \to \mathbb{P}^{\log(n)}$ (example only)

 $\bullet$  e.g. standard encoding for b  $\in \, \mathbb{B}^n$ 

#### individual:

010	101	111	000	110	001	101	100	← genotype
0	1	2	3	4	5	6	7	← index

consider index and associated genotype entry as unit / record / struct; sort units with respect to genotype value, old indices yield permutation:

3 5 0 7 1 6 4 2 ← old index	ſ	000	001	010	100	101	101	110	111	←— genotype
		٠.	5	0	7	1	6	4	2	← old index

= permutation

### **Design of Evolutionary Algorithms**

Lecture 06

ad 1a) genotype-phenotype mapping

typically required: strong causality

- → small changes in individual leads to small changes in fitness
- → small changes in genotype should lead to small changes in phenotype

but: how to find a genotype-phenotype mapping with that property?

#### necessary conditions:

technische universität

- 1) g:  $\mathbb{B}^n \to X$  can be computed efficiently (otherwise it is senseless)
- 2) g:  $\mathbb{B}^n \to X$  is surjective (otherwise we might miss the optimal solution)
- 3) g:  $\mathbb{B}^n \to X$  preserves closeness (otherwise strong causality endangered)

Let  $d(\cdot, \cdot)$  be a metric on  $\mathbb{B}^n$  and  $d_x(\cdot, \cdot)$  be a metric on X.

 $\forall x, y, z \in \mathbb{B}^n : d(x, y) \le d(x, z) \Rightarrow d_X(g(x), g(y)) \le d_X(g(x), g(z))$ 

#### **Design of Evolutionary Algorithms**

Lecture 06

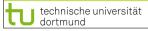
ad 1b) use "most natural" representation

typically required: strong causality

- → small changes in individual leads to small changes in fitness
- → need variation operators that obey that requirement

**but**: how to find variation operators with that property?

 $\Rightarrow$  need design guidelines ...



G. Rudolph: Computational Intelligence • Winter Term 2021/22

#### Design of Evolutionary Algorithms

Lecture 06

#### ad 2) design guidelines for variation operators

#### a) reachability

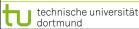
every  $x \in X$  should be reachable from arbitrary  $x_0 \in X$  after finite number of repeated variations with positive probability bounded from 0

#### b) unbiasedness

unless having gathered knowledge about problem variation operator should not favor particular subsets of solutions ⇒ formally: maximum entropy principle

#### c) control

variation operator should have parameters affecting shape of distributions; known from theory: weaken variation strength when approaching optimum



G. Rudolph: Computational Intelligence • Winter Term 2021/22

10

#### **Design of Evolutionary Algorithms**

Lecture 06

#### ad 2) design guidelines for variation operators in practice

binary search space  $X = \mathbb{B}^n$ 

variation by k-point or uniform crossover and subsequent mutation

#### a) reachability:

regardless of the output of crossover we can move from  $x \in B^n$  to  $y \in B^n$  in 1 step with probability

$$p(x,y) = p_m^{H(x,y)} (1 - p_m)^{n - H(x,y)} > 0$$

where H(x,y) is Hamming distance between x and y.

Since  $\min\{p(x,y): x,y \in \mathbb{B}^n\} = \delta > 0$  we are done.

#### **Design of Evolutionary Algorithms**

Lecture 06

#### b) unbiasedness

don't prefer any direction or subset of points without reason

⇒ use maximum entropy distribution for sampling!

#### properties:

technische universität

- distributes probability mass as uniform as possible
- additional knowledge can be included as constraints:
- $\rightarrow$  under given constraints sample as uniform as possible

#### **Design of Evolutionary Algorithms**

Lecture 06

Formally:

#### **Definition:**

Let X be discrete random variable (r.v.) with  $p_k = P\{X = x_k\}$  for some index set K. The quantity

 $H(X) = -\sum_{k \in K} p_k \log p_k$ 

is called the *entropy of the distribution* of X. If X is a continuous r.v. with p.d.f.  $f_{x}(\cdot)$  then the entropy is given by

$$H(X) = -\int_{-\infty}^{\infty} f_X(x) \log f_X(x) dx$$

The distribution of a random variable X for which H(X) is maximal is termed a maximum entropy distribution.



#### **Excursion: Maximum Entropy Distributions**

Lecture 06

#### Knowledge available:

Discrete distribution with support  $\{x_1, x_2, \dots x_n\}$  with  $x_1 < x_2 < \dots x_n < \infty$ 

$$p_k = P\{X = x_k\}$$

⇒ leads to nonlinear constrained optimization problem:

$$-\sum_{k=1}^n p_k \log p_k 
ightarrow \max!$$
 s.t.  $\sum_{k=1}^n p_k = 1$ 

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p, a) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right)$$



G. Rudolph: Computational Intelligence • Winter Term 2021/22

#### **Excursion: Maximum Entropy Distributions**

# Lecture 06

$$L(p,a) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right)$$

partial derivatives:

U technische universität dortmund

$$\frac{\partial L(p,a)}{\partial p_k} = -1 - \log p_k + a \stackrel{!}{=} 0 \qquad \Rightarrow p_k \stackrel{!}{=} e^{a-1}$$
 
$$\frac{\partial L(p,a)}{\partial a} = \sum_{k=1}^n p_k - 1 \stackrel{!}{=} 0$$
 
$$p_k = \frac{1}{n}$$
 uniform distribution 
$$\frac{1}{n} \sum_{k=1}^n p_k = \sum_{k=1}^n e^{a-1} = n e^{a-1} \stackrel{!}{=} 1 \qquad \Leftrightarrow \qquad e^{a-1} = \frac{1}{n}$$

#### **Excursion: Maximum Entropy Distributions**

#### Lecture 06

#### Knowledge available:

technische universität dortmund

Discrete distribution with support  $\{1, 2, ..., n\}$  with  $p_k = P\{X = k\}$  and E[X] = v

⇒ leads to nonlinear constrained optimization problem:

$$-\sum\limits_{k=1}^n p_k \log p_k \longrightarrow ext{max!}$$
 s.t.  $\sum\limits_{k=1}^n p_k = 1$  and  $\sum\limits_{k=1}^n k \, p_k = 
u$ 

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p, a, b) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right) + b \left(\sum_{k=1}^{n} k \cdot p_k - \nu\right)$$

#### Lecture 06

$$L(p, a, b) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right) + b \left(\sum_{k=1}^{n} k \cdot p_k - \nu\right)$$

partial derivatives:

$$\frac{\partial L(p,a,b)}{\partial p_k} = -1 - \log p_k + a + b k \stackrel{!}{=} 0 \qquad \Rightarrow p_k = e^{a-1+b k}$$

$$\frac{\partial L(p,a,b)}{\partial a} = \sum_{k=1}^{n} p_k - 1 \stackrel{!}{=} 0$$

$$\frac{\partial L(p,a,b)}{\partial b} \stackrel{(*)}{=} \sum_{k=1}^{n} k p_k - \nu \stackrel{!}{=} 0 \qquad \sum_{k=1}^{n} p_k = e^{a-1} \sum_{k=1}^{n} (e^b)^k \stackrel{!}{=} 1$$

(continued on next slide)

technische universität dortmund G. Rudolph: Computational Intelligence • Winter Term 2021/22

#### **Excursion: Maximum Entropy Distributions**

Lecture 06

$$\Rightarrow e^{a-1} = \frac{1}{\sum_{k=1}^{n} (e^b)^k} \qquad \Rightarrow p_k = e^{a-1+bk} = \frac{(e^b)^k}{\sum_{i=1}^{n} (e^b)^i}$$

$$\Rightarrow$$
 discrete Boltzmann distribution  $p_k = rac{q^k}{\sum\limits_{i=1}^n q^i}$   $(q=e^b)$ 

 $\Rightarrow$  value of q depends on v via third condition: ( $\star$ )

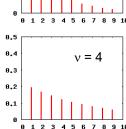
$$\sum_{k=1}^{n} k p_{k} = \frac{\sum_{k=1}^{n} k q^{k}}{\sum_{i=1}^{n} q^{i}} = \frac{1 - (n+1) q^{n} + n q^{n+1}}{(1-q)(1-q^{n})} \stackrel{!}{=} \nu$$

U technische universität dortmund G. Rudolph: Computational Intelligence • Winter Term 2021/22

# **Excursion: Maximum Entropy Distributions**

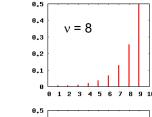
#### θ.5 θ.4 θ.3 θ.2 θ.1 θ 1 2 3 4 5 6 7 8 9 1θ

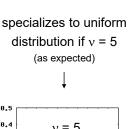
# 0 0 1 2 3 4 5 6 7 8 9 10 0 1 2 3 4 5 6 7 8 9 10 0 5 0 1 2 3 4 5 6 7 8 9 10



technische universität

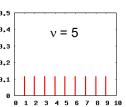
# Distributions Lecture 06



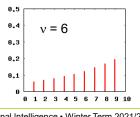


**Boltzmann distribution** 

(n = 9)



# 



#### G. Rudolph: Computational Intelligence • Winter Term 2021/22

#### **Excursion: Maximum Entropy Distributions**

#### Lecture 06

#### Knowledge available:

Discrete distribution with support { 1, 2, ..., n } with E[X] = v and V[X] =  $\eta^2$ 

 $\Rightarrow$  leads to nonlinear constrained optimization problem:

$$-\sum_{k=1}^n p_k \log p_k \to \max!$$
 s.t.  $\sum_{k=1}^n p_k = 1$  and  $\sum_{k=1}^n k p_k = \nu$  and  $\sum_{k=1}^n (k-\nu)^2 p_k = \eta^2$ 

<u>solution:</u> in principle, via Lagrange (find stationary point of Lagrangian function)

but very complicated analytically, if possible at all

⇒ consider special cases only

**note:** constraints are linear equations in p<sub>k</sub>

Lecture 06

Special case: n = 3 and E[X] = 2 and  $V[X] = n^2$ 

Linear constraints uniquely determine distribution:

I. 
$$p_1 + p_2 + p_3 = 1$$
  
II.  $p_1 + 2p_2 + 3p_3 = 2$   
III.  $p_1 + 0 + p_3 = \eta^2$   
III.  $p_1 + 0 + p_3 = \eta^2$   
III.  $p_2 + 2p_3 = 1$   
 $p_2 + 2p_3 = 1$   
 $p_3 = \frac{\eta^2}{2}$ 

$$\Rightarrow p = \left(\frac{\eta^2}{2}, 1 - \eta^2, \frac{\eta^2}{2}\right) \qquad \begin{array}{c} \eta^2 = \frac{1}{4} & \eta^2 = \frac{2}{3} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta_2 \cdot \theta} \\ \frac{\theta_1 \cdot \theta}{\theta_2 \cdot \theta} & \frac{\theta_2 \cdot \theta}{\theta} \\ \frac{\theta_1 \cdot \theta}{\theta} & \frac{\theta}{\theta} \\ \frac{\theta}{\theta} & \frac{\theta}{\theta} \\ \frac{\theta}{$$

technische universität dortmund

G. Rudolph: Computational Intelligence • Winter Term 2021/22

Lecture 06

#### **Excursion: Maximum Entropy Distributions**

$$L(p,a,b) = -\sum_{k=0}^{\infty} p_k \log p_k + a \left(\sum_{k=0}^{\infty} p_k - 1\right) + b \left(\sum_{k=0}^{\infty} k \cdot p_k - \nu\right)$$

partial derivatives:

$$\frac{\partial L(p,a,b)}{\partial p_k} = -1 - \log p_k + a + b k \stackrel{!}{=} 0 \qquad \Rightarrow p_k = e^{a-1+b k}$$

$$\frac{\partial L(p,a,b)}{\partial a} = \sum_{k=0}^{\infty} p_k - 1 \stackrel{!}{=} 0$$

$$\frac{\partial L(p,a,b)}{\partial a} = \sum_{k=0}^{\infty} p_k - 1 \stackrel{!}{=} 0$$

$$\frac{\partial L(p,a,b)}{\partial b} \stackrel{(*)}{=} \sum_{k=0}^{\infty} k p_k - \nu \stackrel{!}{=} 0$$

$$\sum_{k=0}^{\infty} p_k = e^{a-1} \sum_{k=0}^{\infty} (e^b)^k \stackrel{!}{=} 1$$

(continued on next slide)

#### **Excursion: Maximum Entropy Distributions**

Lecture 06

#### Knowledge available:

Discrete distribution with unbounded support  $\{0, 1, 2, ...\}$  and E[X] = v

⇒ leads to infinite-dimensional nonlinear constrained optimization problem:

$$-\sum_{k=0}^{\infty}p_k\log p_k \to \max!$$
 s.t. 
$$\sum_{k=0}^{\infty}p_k = 1 \quad \text{and} \quad \sum_{k=0}^{\infty}k\,p_k = \nu$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p,a,b) = -\sum_{k=0}^{\infty} p_k \log p_k + a \left(\sum_{k=0}^{\infty} p_k - 1\right) + b \left(\sum_{k=0}^{\infty} k \cdot p_k - \nu\right)$$

technische universität

G. Rudolph: Computational Intelligence • Winter Term 2021/22

#### **Excursion: Maximum Entropy Distributions**

Lecture 06

$$\Rightarrow e^{a-1} = \frac{1}{\sum_{k=0}^{\infty} (e^b)^k} \qquad \Rightarrow p_k = e^{a-1+bk} = \frac{(e^b)^k}{\sum_{i=0}^{\infty} (e^b)^i}$$

set 
$$q=e^b$$
 and insists that  $q<1$   $\Rightarrow$   $\sum_{k=0}^{\infty}q^k$   $=$   $\frac{1}{1-q}$  insert

$$p_k = (1-q)\,q^k$$
 for  $k=0,1,2,\ldots$  geometrical distribution

it remains to specify q; to proceed recall that  $\sum_{k=0}^{\infty} k \, q^k \, = \, \frac{q}{(1-a)^2}$ 

Lecture 06

value of q depends on v via third condition: (\*)

$$\sum_{k=0}^{\infty} k \, p_k \, = \, \frac{\sum_{k=0}^{\infty} k \, q^k}{\sum_{i=0}^{\infty} q^i} \, = \, \frac{q}{1-q} \, \stackrel{!}{=} \, \nu$$

$$\Rightarrow q = \frac{\nu}{\nu + 1} = 1 - \frac{1}{\nu + 1}$$

$$\Rightarrow p_k = \frac{1}{\nu+1} \left( 1 - \frac{1}{\nu+1} \right)^k$$



G. Rudolph: Computational Intelligence • Winter Term 2021/22

#### Lecture 06 **Excursion: Maximum Entropy Distributions**

#### Overview:

support { 1, 2, ..., n } ⇒ discrete uniform distribution

and require  $E[X] = \theta$ ⇒ Boltzmann distribution

⇒ N.N. (**not** Binomial distribution) and require  $V[X] = \eta^2$ 

⇒ not defined! support N

and require  $E[X] = \theta$ ⇒ *geometrical* distribution

and require  $V[X] = \eta^2$  $\Rightarrow$  ?

support  $\mathbb{Z}$  $\Rightarrow$  not defined!

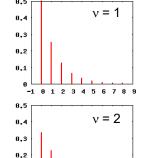
and require  $E[|X|] = \theta$ ⇒ bi-geometrical distribution (discrete Laplace distr.)

and require  $E[|X|^2] = \eta^2$ ⇒ N.N. (discrete Gaussian distr.)

# technische universität

# **Excursion: Maximum Entropy Distributions**

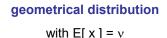


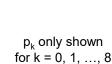


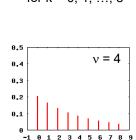
technische universität

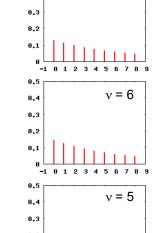
dortmund

v = 3









v = 7

G. Rudolph: Computational Intelligence • Winter Term 2021/22

#### **Excursion: Maximum Entropy Distributions**

Lecture 06

support [a,b]  $\subset \mathbb{R}$ ⇒ uniform distribution

support  $\mathbb{R}^+$  with  $E[X] = \theta \implies$  Exponential distribution

support R

0.3

with E[X] =  $\theta$ , V[X] =  $\eta^2$   $\Rightarrow$  normal / Gaussian distribution N( $\theta$ ,  $\eta^2$ )

support Rn

with  $E[X] = \theta$ 

technische universität

and Cov[X] = C $\Rightarrow$  multinormal distribution N( $\theta$ , C)

expectation vector  $\in \mathbb{R}^n$ 

covariance matrix  $\in \mathbb{R}^{n,n}$ positive definite:

 $\forall x \neq 0 : x'Cx > 0$ 

Lecture 06

for permutation distributions?

ightarrow uniform distribution on all possible permutations

```
 \begin{array}{l} \text{set } \mathbf{v}[j] = j \text{ for } j = 1,\ 2,\ \ldots,\ n \\ \\ \text{for } i = n \text{ to } 1 \text{ step } -1 \\ \\ \text{draw } k \text{ uniformly at random from } \{\ 1,\ 2,\ \ldots,\ i\ \} \\ \\ \text{swap } \mathbf{v}[i] \text{ and } \mathbf{v}[k] \\ \\ \text{endfor} \\ \end{array} \right) \\ \begin{array}{l} \text{generates permutation uniformly at random in } \\ \\ \Theta(n) \text{ time} \\ \end{array}
```

#### Guideline:

Only if you know something about the problem a priori or

if you have learnt something about the problem during the search

⇒ include that knowledge in search / mutation distribution (via constraints!)



G. Rudolph: Computational Intelligence • Winter Term 2021/22

29

