

# **Computational Intelligence**

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- Design of Evolutionary Algorithms
  - Design Guidelines
  - Genotype-Phenotype Mapping
  - Maximum Entropy Distributions

#### Three tasks:

- 1. Choice of an appropriate problem representation.
- 2. Choice / design of variation operators acting in problem representation.
- 3. Choice of strategy parameters (includes initialization).

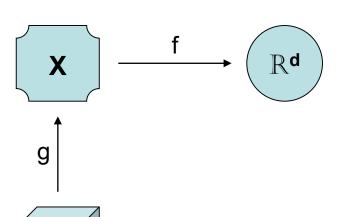
#### ad 1) different "schools":

- (a) operate on binary representation and define genotype/phenotype mapping
  - + can use standard algorithm
  - mapping may induce unintentional bias in search
- (b) no doctrine: use "most natural" representation
  - must design variation operators for specific representation
  - + if design done properly then no bias in search

ad 1a) genotype-phenotype mapping

original problem  $f: X \to \mathbb{R}^d$ 

scenario: no standard algorithm for search space X available



- standard EA performs variation on binary strings b  $\in \mathbb{B}^n$
- fitness evaluation of individual b via  $(f \circ g)(b) = f(g(b))$ where g:  $\mathbb{B}^n \to X$  is genotype-phenotype mapping
- selection operation independent from representation

 $\mathbb{B}$ n

## **Genotype-Phenotype-Mapping** $\mathbb{B}^n \to [L, R] \subset \mathbb{R}$

• Standard encoding for  $b \in \mathbb{B}^n$ 

$$x = L + \frac{R - L}{2^n - 1} \sum_{i=0}^{n-1} b_{n-i} 2^i$$

→ Problem: *hamming cliffs* 

000	001	010	011	100	101	110	111		
0	1	1 2		3 4		6	7		
0 1 2 3 4 5 6 7  1 Bit 2 Bit 1 Bit 3 Bit 1 Bit 2 Bit 1 Bit									

denotype denotype

$$L = 0, R = 7$$
  
 $n = 3$ 

## **Genotype-Phenotype-Mapping** $\mathbb{B}^n \to [L, R] \subset \mathbb{R}$

• Gray encoding for  $b \in \mathbb{B}^n$ 

Let 
$$a \in \mathbb{B}^n$$
 standard encoded. Then  $b_i = \left\{ \begin{array}{l} a_i, & \text{if } i = 1 \\ a_{i-1} \oplus a_i, & \text{if } i > 1 \end{array} \right. \oplus = XOR$ 

000	001	011	010	110	111	101	100	← genotype
0	1	2	3	4	5	6	7	← phenotype

OK, no hamming cliffs any longer ...

- ⇒ small changes in phenotype "lead to" small changes in genotype since we consider evolution in terms of Darwin (not Lamarck):
- ⇒ small changes in genotype lead to small changes in phenotype!

**but:** 1-Bit-change:  $000 \rightarrow 100 \Rightarrow \odot$ 

**Genotype-Phenotype-Mapping**  $\mathbb{B}^n \to \mathbb{P}^{\log(n)}$  (example only)

ullet e.g. standard encoding for  $b \in \mathbb{B}^n$ 

#### individual:

010	101	111	000	110	001	101	100	← genotype
0	1	2	3	4	5	6	7	← index

consider index and associated genotype entry as unit / record / struct; sort units with respect to genotype value, old indices yield permutation:

000	001	010	100	101	101	110	111	← genotype
3	5	0	7	1	6	4	2	← old index

= permutation

ad 1a) genotype-phenotype mapping

typically required: strong causality

- → small changes in individual leads to small changes in fitness
- → small changes in genotype should lead to small changes in phenotype

but: how to find a genotype-phenotype mapping with that property?

#### necessary conditions:

- 1) g:  $\mathbb{B}^n \to X$  can be computed efficiently (otherwise it is senseless)
- 2) g:  $\mathbb{B}^n \to X$  is surjective (otherwise we might miss the optimal solution)
- 3) g:  $\mathbb{B}^n \to X$  preserves closeness (otherwise strong causality endangered)

Let  $d(\cdot, \cdot)$  be a metric on  $\mathbb{B}^n$  and  $d_X(\cdot, \cdot)$  be a metric on X.

 $\forall x, y, z \in \mathbb{B}^n : d(x, y) \le d(x, z) \Rightarrow d_X(g(x), g(y)) \le d_X(g(x), g(z))$ 

# **Design of Evolutionary Algorithms**

Lecture 06

ad 1b) use "most natural" representation

typically required: strong causality

- → small changes in individual leads to small changes in fitness
- → need variation operators that obey that requirement

**but**: how to find variation operators with that property?

⇒ need design guidelines ...

## ad 2) design guidelines for variation operators

#### a) reachability

every  $x \in X$  should be reachable from arbitrary  $x_0 \in X$  after finite number of repeated variations with positive probability bounded from 0

#### b) unbiasedness

unless having gathered knowledge about problem variation operator should not favor particular subsets of solutions ⇒ formally: maximum entropy principle

#### c) control

variation operator should have parameters affecting shape of distributions; known from theory: weaken variation strength when approaching optimum

## ad 2) design guidelines for variation operators in practice

binary search space  $X = \mathbb{B}^n$ 

variation by k-point or uniform crossover and subsequent mutation

#### a) reachability:

regardless of the output of crossover we can move from  $x \in \mathbb{B}^n$  to  $y \in \mathbb{B}^n$  in 1 step with probability

$$p(x,y) = p_m^{H(x,y)} (1 - p_m)^{n - H(x,y)} > 0$$

where H(x,y) is Hamming distance between x and y.

Since  $\min\{p(x,y): x,y \in \mathbb{B}^n\} = \delta > 0$  we are done.

#### b) *unbiasedness*

don't prefer any direction or subset of points without reason

⇒ use maximum entropy distribution for sampling!

#### properties:

- distributes probability mass as uniform as possible
- additional knowledge can be included as constraints:
  - → under given constraints sample as uniform as possible

#### Formally:

#### **Definition:**

Let X be discrete random variable (r.v.) with  $p_k = P\{X = x_k\}$  for some index set K. The quantity

$$H(X) = -\sum_{k \in K} p_k \log p_k$$

is called the **entropy of the distribution** of X. If X is a continuous r.v. with p.d.f.  $f_X(\cdot)$  then the entropy is given by

$$H(X) = -\int_{-\infty}^{\infty} f_X(x) \log f_X(x) dx$$

The distribution of a random variable X for which H(X) is maximal is termed a *maximum entropy distribution*.

#### Knowledge available:

Discrete distribution with support {  $x_1, x_2, ... x_n$  } with  $x_1 < x_2 < ... x_n < \infty$   $p_k = P\{X = x_k\}$ 

⇒ leads to nonlinear constrained optimization problem:

$$-\sum_{k=1}^n p_k \log p_k 
ightarrow \max!$$
 s.t.  $\sum_{k=1}^n p_k = 1$ 

solution: via Lagrange (find stationary point of Lagrangian function)

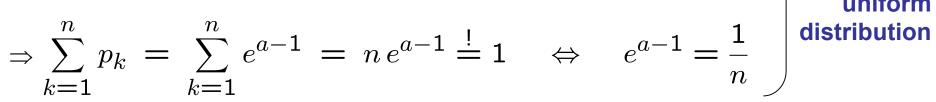
$$L(p,a) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right)$$

$$L(p,a) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right)$$

partial derivatives:

$$\frac{\partial L(p,a)}{\partial p_k} = -1 - \log p_k + a \stackrel{!}{=} 0 \qquad \Rightarrow p_k \stackrel{!}{=} e^{a-1}$$

$$\frac{\partial L(p,a)}{\partial a} = \sum_{k=1}^{n} p_k - 1 \stackrel{!}{=} 0$$





#### Knowledge available:

Discrete distribution with support  $\{1, 2, ..., n\}$  with  $p_k = P\{X = k\}$  and E[X] = v

⇒ leads to nonlinear constrained optimization problem:

$$-\sum_{k=1}^n p_k \log p_k \to \max!$$
 s.t. 
$$\sum_{k=1}^n p_k = 1 \qquad \text{and} \qquad \sum_{k=1}^n k \, p_k = \nu$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p, a, b) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right) + b \left(\sum_{k=1}^{n} k \cdot p_k - \nu\right)$$

$$L(p, a, b) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right) + b \left(\sum_{k=1}^{n} k \cdot p_k - \nu\right)$$

partial derivatives:

$$\frac{\partial L(p,a,b)}{\partial p_k} = -1 - \log p_k + a + b \cdot k \stackrel{!}{=} 0 \qquad \Rightarrow p_k = e^{a-1+b \cdot k}$$

$$\frac{\partial L(p,a,b)}{\partial a} = \sum_{k=1}^{n} p_k - 1 \stackrel{!}{=} 0$$

$$\frac{\partial L(p,a,b)}{\partial b} \stackrel{(*)}{=} \sum_{k=1}^{n} k p_k - \nu \stackrel{!}{=} 0 \qquad \sum_{k=1}^{n} p_k = e^{a-1} \sum_{k=1}^{n} (e^b)^k \stackrel{!}{=} 1$$

(continued on next slide)

# **Excursion: Maximum Entropy Distributions**

#### Lecture 06

$$\Rightarrow e^{a-1} = \frac{1}{\sum_{k=1}^{n} (e^b)^k} \Rightarrow p_k = e^{a-1+bk} = \frac{(e^b)^k}{\sum_{i=1}^{n} (e^b)^i}$$

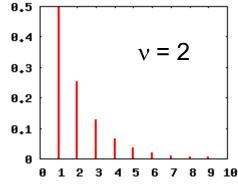
$$\Rightarrow$$
 discrete Boltzmann distribution  $p_k = \frac{q^k}{\sum\limits_{i=1}^n q^i}$   $(q = e^b)$ 

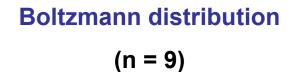
 $\Rightarrow$  value of q depends on v via third condition: (\*)

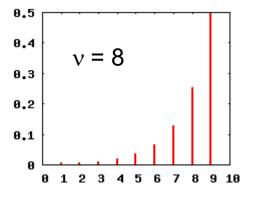
$$\sum_{k=1}^{n} k p_k = \frac{\sum_{k=1}^{n} k q^k}{\sum_{i=1}^{n} q^i} = \frac{1 - (n+1) q^n + n q^{n+1}}{(1-q)(1-q^n)} \stackrel{!}{=} \nu$$

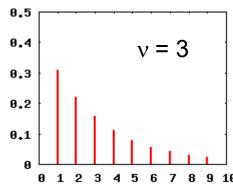
# **Excursion: Maximum Entropy Distributions**

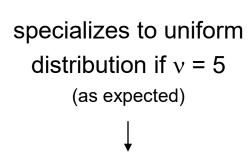
## **Lecture 06**

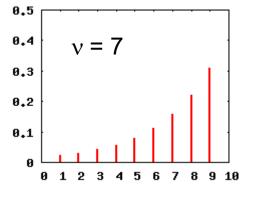


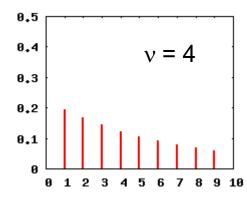


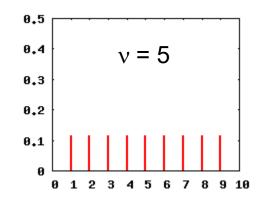


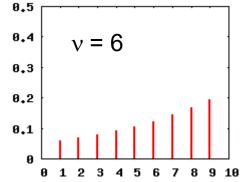












#### **Knowledge available:**

Discrete distribution with support { 1, 2, ..., n } with E[X] = v and  $V[X] = \eta^2$ 

⇒ leads to nonlinear constrained optimization problem:

$$-\sum_{k=1}^n p_k \log p_k \longrightarrow \max!$$

s.t. 
$$\sum_{k=1}^n p_k = 1 \quad \text{and} \quad \sum_{k=1}^n k \, p_k = \nu \quad \text{and} \quad \sum_{k=1}^n (k-\nu)^2 \, p_k = \eta^2$$

solution: in principle, via Lagrange (find stationary point of Lagrangian function)

but very complicated analytically, if possible at all

⇒ consider special cases only

**note:** constraints are linear equations in p<sub>k</sub>

Special case: n = 3 and E[X] = 2 and  $V[X] = \eta^2$ 

Linear constraints uniquely determine distribution:

I. 
$$p_1 + p_2 + p_3 = 1$$
  
II.  $p_1 + 2p_2 + 3p_3 = 2$   
III.  $p_1 + 0 + p_3 = \eta^2$   $p_1 = \frac{\eta^2}{2}$   
IIII.  $p_2 + 2p_3 = 1$   $p_3 = \frac{\eta^2}{2}$   $p_3 = \frac{\eta^2}{2}$ 

$$\Rightarrow p = \left(\frac{\eta^2}{2}, 1 - \eta^2, \frac{\eta^2}{2}\right) \qquad \begin{array}{c} \eta^2 = \frac{1}{4} & \eta^2 = \frac{2}{3} & \eta^2 = \frac{4}{5} \\ \frac{\theta \cdot \theta}{\theta \cdot \theta} & \frac{\theta \cdot \theta}{\theta} & \frac{\theta}{\theta} & \frac{\theta}{$$

#### Knowledge available:

Discrete distribution with unbounded support  $\{0, 1, 2, ...\}$  and E[X] = v

⇒ leads to infinite-dimensional nonlinear constrained optimization problem:

$$-\sum_{k=0}^\infty p_k \log p_k \to \max!$$
 s.t. 
$$\sum_{k=0}^\infty p_k = 1 \qquad \text{and} \qquad \sum_{k=0}^\infty k \, p_k = \nu$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p, a, b) = -\sum_{k=0}^{\infty} p_k \log p_k + a \left(\sum_{k=0}^{\infty} p_k - 1\right) + b \left(\sum_{k=0}^{\infty} k \cdot p_k - \nu\right)$$

# **Excursion: Maximum Entropy Distributions**

## Lecture 06

$$L(p, a, b) = -\sum_{k=0}^{\infty} p_k \log p_k + a \left(\sum_{k=0}^{\infty} p_k - 1\right) + b \left(\sum_{k=0}^{\infty} k \cdot p_k - \nu\right)$$

partial derivatives:

$$\frac{\partial L(p,a,b)}{\partial p_k} = -1 - \log p_k + a + b \cdot k \stackrel{!}{=} 0 \qquad \Rightarrow p_k = e^{a-1+b \cdot k}$$

$$\frac{\partial L(p,a,b)}{\partial a} = \sum_{k=0}^{\infty} p_k - 1 \stackrel{!}{=} 0$$

$$\frac{\partial L(p,a,b)}{\partial b} \stackrel{(*)}{=} \sum_{k=0}^{\infty} k p_k - \nu \stackrel{!}{=} 0 \qquad \sum_{k=0}^{\infty} p_k = e^{a-1} \sum_{k=0}^{\infty} (e^b)^k \stackrel{!}{=} 1$$

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# **Excursion: Maximum Entropy Distributions**

#### **Lecture 06**

$$\Rightarrow e^{a-1} = \frac{1}{\sum_{k=0}^{\infty} (e^b)^k}$$

$$\Rightarrow p_k = e^{a-1+bk} = \frac{(e^b)^k}{\sum_{i=0}^{\infty} (e^b)^i}$$

set 
$$q=e^b$$
 and insists that  $q<1$   $\Rightarrow$   $\sum_{k=0}^{\infty}q^k$   $=$   $\frac{1}{1-q}$  insert

$$\Rightarrow p_k = (1-q) \, q^k$$
 for  $k = 0, 1, 2, \ldots$  geometrical distribution

it remains to specify q; to proceed recall that 
$$\sum_{k=0}^{\infty} k \, q^k \, = \, \frac{q}{(1-q)^2}$$

 $\Rightarrow$  value of q depends on v via third condition: (\*)

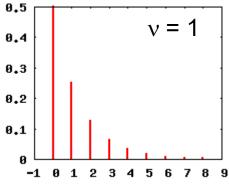
$$\sum_{k=0}^{\infty} k p_k = \frac{\sum_{k=0}^{\infty} k q^k}{\sum_{i=0}^{\infty} q^i} = \frac{q}{1-q} \stackrel{!}{=} \nu$$

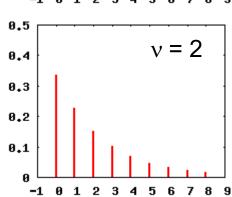
$$\Rightarrow q = \frac{\nu}{\nu + 1} = 1 - \frac{1}{\nu + 1}$$

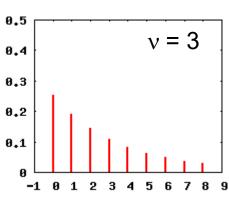
$$\Rightarrow p_k = \frac{1}{\nu+1} \left( 1 - \frac{1}{\nu+1} \right)^k$$

# **Excursion: Maximum Entropy Distributions**

## **Lecture 06**



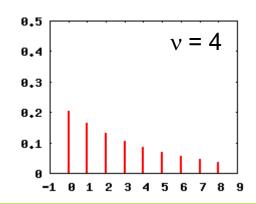


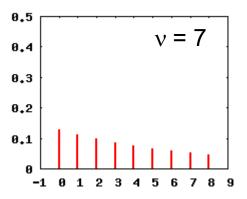


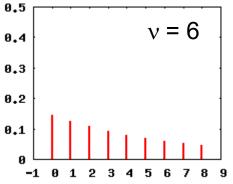
# geometrical distribution

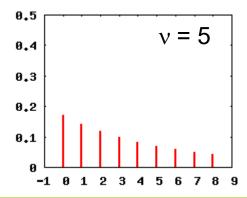
with E[x] = 
$$v$$

$$p_k$$
 only shown for  $k = 0, 1, ..., 8$ 









#### Overview:

support  $\{1, 2, ..., n\}$   $\Rightarrow$  discrete uniform distribution

and require  $E[X] = \theta$   $\Rightarrow$  *Boltzmann* distribution

and require  $V[X] = \eta^2$   $\Rightarrow$  N.N. (**not** Binomial distribution)

support  $\mathbb{N}$   $\Rightarrow$  not defined!

and require  $E[X] = \theta$   $\Rightarrow$  *geometrical* distribution

and require  $V[X] = \eta^2 \Rightarrow ?$ 

support  $\mathbb{Z}$   $\Rightarrow$  not defined!

and require  $E[|X|] = \theta$   $\Rightarrow$  *bi-geometrical* distribution (*discrete Laplace* distr.)

and require  $E[|X|^2] = \eta^2 \Rightarrow N.N.$  (discrete Gaussian distr.)

support [a,b] 
$$\subset \mathbb{R}$$

⇒ uniform distribution

support 
$$\mathbb{R}^+$$
 with  $\mathsf{E}[\mathsf{X}] = \theta$ 

support  $\mathbb{R}^+$  with  $E[X] = \theta \implies Exponential distribution$ 

support  $\mathbb{R}$ 

with E[X] = 
$$\theta$$
, V[X] =  $\eta^2$ 

with E[X] =  $\theta$ , V[X] =  $\eta^2$   $\Rightarrow$  normal / Gaussian distribution N( $\theta$ ,  $\eta^2$ )

support  $\mathbb{R}^n$ 

with  $E[X] = \theta$ 

and Cov[X] = C

 $\Rightarrow$  multinormal distribution N( $\theta$ , C)

expectation vector  $\in \mathbb{R}^n$ 

covariance matrix  $\in \mathbb{R}^{n,n}$ 

positive definite:

 $\forall x \neq 0 : x'Cx > 0$ 

for permutation distributions?

→ uniform distribution on all possible permutations

#### **Guideline:**

Only if you know something about the problem *a priori* or if you have learnt something about the problem *during the search* 

⇒ include that knowledge in search / mutation distribution (via constraints!)